Table S1: Mathematical definitions of complex network statistics used in the present study

| Ref. | Nomenclature | Description of use |
| :---: | :---: | :---: |
| $N / A$ | $n$ | Number of nodes (All networks were calculated on the basis of 59 possible nodes, corresponding to the 59 recordable electrodes). The set of nodes in the network is denoted $N$ (this is all nodes for which links were found), and the number of nodes $n$. |
| $N / A$ | $m$ | Number of links (edges) in the network (all networks were calculated on the basis of 59 * 58 (3422) possible links (a node was not allowed to have a link to itself). Networks compared were directed symmetrical; all links were counted in both directions to avoid ambiguity with directed links. The set of links in the network is denoted $M$, and the number of links found $m^{\dagger}$ (each counted twice). |
| N/A | $\xi$ | Edge density of the network (the same as 'cost' used in [8] is the number of edges found compared to the maximum possible number of edges (given the number of nodes in the networks). For the symmetrical networks used herein, this was calculated as: $m / n(n-1)$ |
| [14] | $a_{i j}$ | The connection status of nodes $i, j, a_{i j}=1$ when a link exists between nodes $i, j$, (they are said to be neighbours). As the networks compared were symmetrical, $a_{i j}==a_{j j}$. |
|  | $k_{i}$ | Degree of node $i$ is the number of neighbors of the node, calculated as: $k_{i}=\sum_{j \in N} a_{i j}$ <br> As nodes must have at least 1 link to be considered part of the network, $k_{i}$ is always $>0$ |


| Ref. | Nomenclature | Description of use |
| :---: | :---: | :---: |
| [14] | $d_{i j}$ | Shortest path length (distance) between nodes i and j , calculated as: $d_{i j}=\sum_{a_{u v} \in g i \leftrightarrow j} a_{u v}$ <br> Where $g_{i \leftrightarrow j}$ is the shortest path (geodesic) between $i$ and $j$. For all disconnected pairs $i, j, d_{i j}=\infty$ |
|  | $I_{i}$ | The path length of node $i$ is the sum of the shortest distance to all other nodes (normalized to $n-1$ ), where shortest distance is calculated as: $I_{i}=1 / n \Sigma\left(\Sigma j d_{i j} / n-1\right)$ <br> Where $j$ is for each other node in the network. |
| [16] | $c^{\text {ws }}$, | Following Watts and Strogatz (ws), the clustering coefficient of node $i$ is the number of edges between the neighbours of $i$ divided by the maximum possible number of edges. This was calculated as: $c_{i}=m / k_{i}\left(k_{i}-1\right)$ <br> Where $m$ is the number of edges between the neighbours of $i$. |
| [14] | L | Mean path length, taken over all node pairs between which a path exists (i.e. $d_{i j} \neq \infty$ ), calculated as: $L=1 / n \sum I_{i}$ <br> Where $l_{i}$ is the average distance between node $i$ and all other nodes. |
|  | C | Mean clustering coefficient, taken over all nodes in the network, calculated as: $C=1 / n \sum c_{i}$ |
|  | K | Mean node degree, taken over all nodes in the network, calculated as: $K=1 / n \sum k_{i}$ |


| Ref. | Nomenclature | Description of use |
| :---: | :---: | :---: |
| [14, 64] | E | Following Latora et al. [64], global efficiency is the average inverse shortest path length, taken over all nodes in the network, calculated as: $E=1 / n \Sigma E_{i}=1 / n \Sigma\left(\Sigma j d_{i j}^{-1} / n-1\right)$ |
| [64] | $C^{\text {ws }}$ Lattice | Expected Watts and Strogatz Clustering Coefficient for an equivalent lattice network, calculated as: $C^{w s}=3 / 4(k-2) /(k-1)$ |
| [14] | $C^{W s}{ }_{\text {Rand }}$ | Expected Watts and Strogatz Clustering Coefficient for an equivalent random network (calculated from a population of random networks with the same number of nodes and links: see section on generation of equivalent null hypothesis networks). |
|  | $L_{\text {Rand }}$ | Expected mean shortest path length for an equivalent random network (calculated from a population of random networks with the same number of nodes and links: see section on generation of equivalent null hypothesis networks). |
| [17] | $S^{w s}=$ | Small Worldness. Calculated as: $\left(C^{\text {ws }} / C^{\text {ws }} \text { Rand }\right) /\left(L / L_{\text {Rand }}\right)$ <br> Using global efficiency instead mean shortest path length [8]: $\left(C^{w s} / C_{\text {Rand }}^{w s}\right) / 1 /\left(E / E_{\text {Rand }}\right)$ |
| N/A | $C S^{\text {ws }}=$ | Conservative Small Worldness. Calculated as: $\left(C^{W s} / C^{\text {Ws }}{ }_{\text {Lattice }}\right) /\left(L / L_{\text {Rand }}\right)$ <br> Using global efficiency instead mean shortest path length [8]: $\left(C^{w s} / C^{\text {Ws }}{ }_{\text {Lattice }}\right) / 1 /\left(E / E_{\text {Rand }}\right)$ |

† NOTE, there is sometimes ambiguity in the notation used for definitions of complex network statistics, with / referring to path length and to number of links, here / is used to refer to path length and $m$ for the number of links.

