**Text S1: maximal meta-module hardness of approximation**

We show here the hardness of the problem of finding a maximal meta-module.

The formal statement of the problem is as follows: we are given two edge-weighted graphs G=(V,E) and G'=(V,E') with the same vertex set V. A meta-module is two disjoint vertex sets S, T in V such that the sum of edge weights of S and T in G is positive and the sum of weights between S and T is positive in G'. The goal is to find a meta-module of maximum cardinality .

**Theorem:** Finding a maximum size meta-module is NP-hard to approximate within any constant factor.

**Proof:** We show that the problem is NP-hard to approximate within a constant factor via a gap preserving reduction from the max-clique problem .Given the input graph G'' = (V'',E'') with *n* nodes for the maximum clique problem, we define the node set for G and G' as where and are copies of V''. For every edge () in E'', let and be the copies of and respectively. In G we set and . In G' we set , , , and . All other edges are scored -. The reduction is clearly polynomial.

Note first that any module in a meta-module cannot contain a negative edge, since the sum of the weights in such a module would be negative. Hence, every module must correspond to a clique in G''. If there is a clique C with at least b nodes in G'', then it will induce a meta-module with at least 2b nodes in G and G', by taking the two copies of C in G and G' as the modules. If a meta-module with at least 2k nodes exists in G and G', then one of its modules has at least k nodes, and such a module corresponds to a clique in G''. In other words, if there is no clique of size a in G'' then there is no meta-module of size 2a in G and G'. Thus, we have shown a gap preserving reduction:

Since the max-clique problem is NP-hard to approximate within a constant factor [[105-107](#_ENREF_105)], we conclude that the maximal meta-module detection is also NP-hard to approximate within a constant factor.